Autonomic and Energy-Efficient Management of Large-Scale Virtualized Data Centers

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Cloud Computing

- On-demand self-service pay-as-you-go resource provisioning
- More and more applications are executed in large data centers



Infrastructure-as-a-Service (laaS) Clouds

- Provide compute capacity in the form of Virtual Machines (VMs)
 - Illusion of a computer running its own operating system
- Server virtualization
 - Multiple VMs on a server
 - Live migration

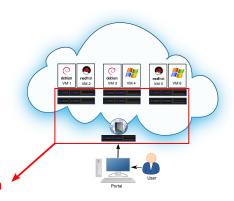


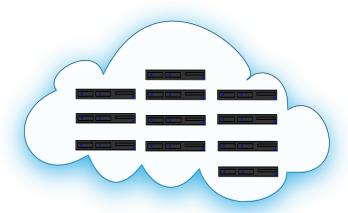
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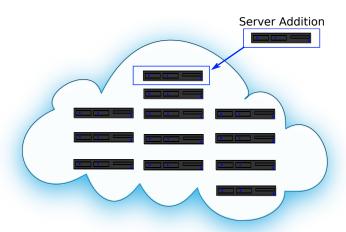
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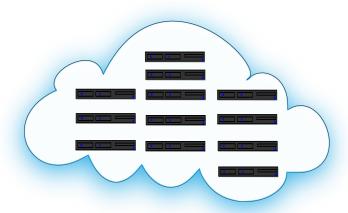
VM management system

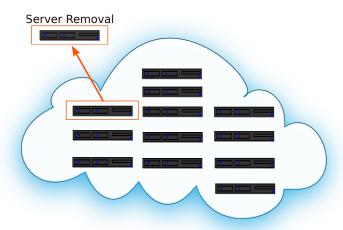
- Controls the servers
- Accepts user requests
- Places VMs on the servers

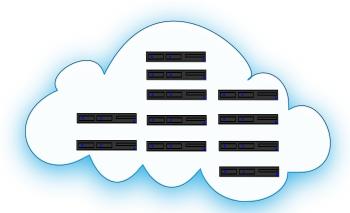














Manual management is impossible

Autonomic laaS cloud management systems are desirable

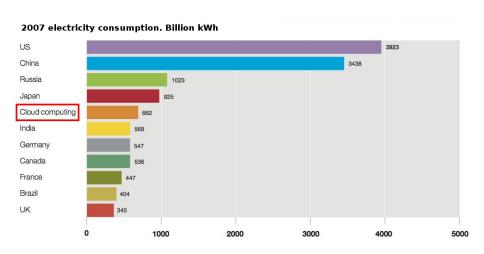
Autonomic System Management

How to achieve autonomic system management in laaS clouds?

- Self-configuration
 - Support for dynamic server addition, removal
- Self-healing
 - Support for automated VM management system services fail-over

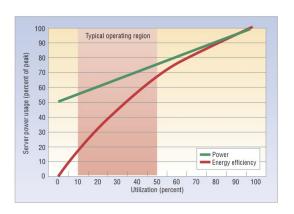
Challenge: Energy Saving

Huge energy amounts in large data centers



Energy Efficiency

- Data centers are rarely fully utilized
 - High fluctuating resource demands → Low utilization (10 to 50%)
- Servers lack power proportionality
 - High idle power consumption
 - Energy efficiency significantly drops under light loads



Energy Saving Approaches

- Slow down the individual server components (e.g. CPU, memory)
 - Becomes less attractive on modern hardware (Le Sueur et al. (2010))

Energy Saving Approaches

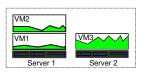
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- Transition parts of the server components into a sleep state
 - Not always easy, as idle time is hard to achieve

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- Slow down the individual server components (e.g. CPU, memory)
 - Becomes less attractive on modern hardware (Le Sueur et al. (2010))
- Transition parts of the server components into a sleep state
 - Not always easy, as idle time is hard to achieve
- Transition entire servers into a sleep state
 - Entering sleep states can yield significant energy savings

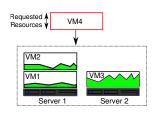
Three methods

 Energy-efficient VM placement



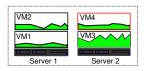
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 Energy-efficient VM placement



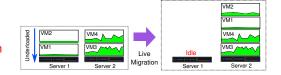
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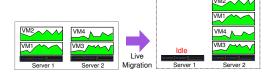
Three methods

- Energy-efficient VM placement
- Server underload detection and mitigation



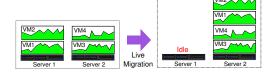
Three methods

- Energy-efficient VM placement
- Server underload detection and mitigation
- Periodic VM consolidation



Three methods

- Energy-efficient VM placement
- Server underload detection and mitigation
- Periodic VM consolidation



Self-optimization for energy efficiency

Objective

Design and implement an autonomic VM management system for large-scale laaS clouds

- Ease of management
- High availability
- Energy efficiency

Contributions

- Snooze: autonomic and energy-efficient VM management system
 - Self-configuring and self-healing VM management system
 - Self-optimizing integrated energy management approach

- Energy-efficient VM management algorithms
 - VM placement
 - VM consolidation

Contributions

- Snooze: autonomic and energy-efficient VM management system
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First Contribution Presented



- Self-configuring and self-healing VM management system
- Self-optimizing integrated energy management approach

System	Architecture	Self-configuration	Self-healing	Evaluation

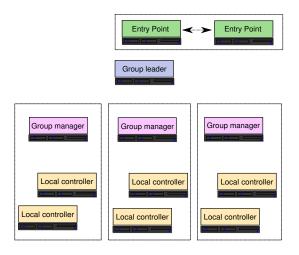
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OpenNebula, OpenStack, Nimbus, Entropy	Centralized	No	No	Real system
CloudStack, VMware DRS	Centralized	No	Yes (Repli- cated servers)	Real system
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Eucalyptus	Static Hierarchy	No	No	Real system

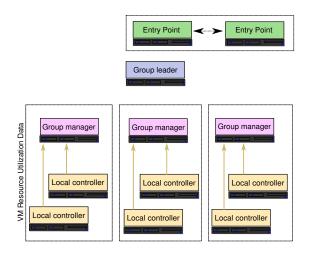
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Snooze	Dynamic Hierarchy	Yes	Yes (No dedi- cated servers)	Real system

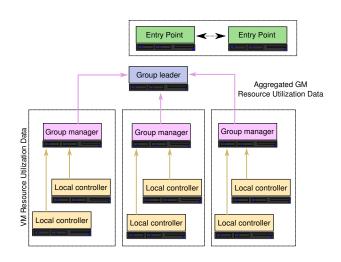
System Architecture



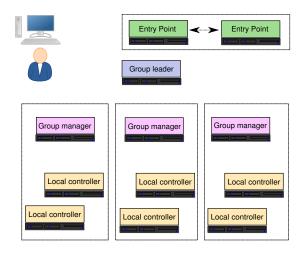
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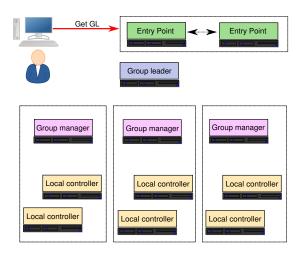
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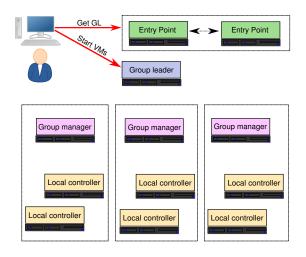
VM Submission Example



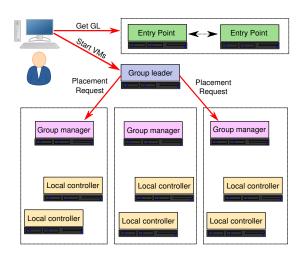
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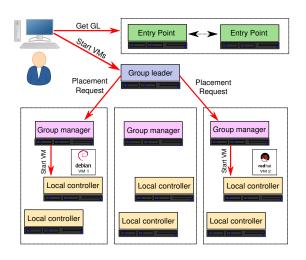
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VM Submission Example



VM Submission Example



Hierarchy Construction and Maintenance

- How to build the hierarchy?
- How to add/remove servers?
- How to deal with server failures?

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Self-configuration and self-healing mechanisms

Hierarchy Construction Protocols

Three steps

- Group leader election
- Group manager join
- Local controller join

- Group leader election algorithm exploiting Apache ZooKeeper
 - Scalable and fault-tolerant coordination framework



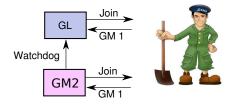
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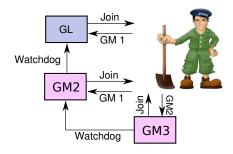
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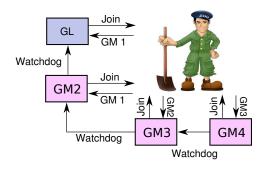
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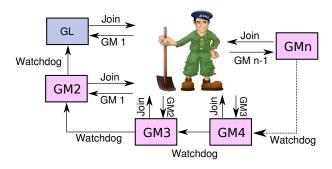
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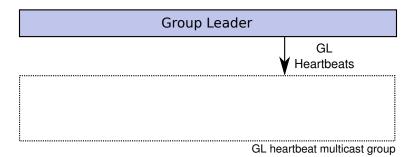
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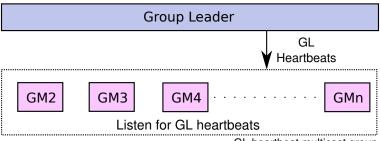
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Group Manager Join

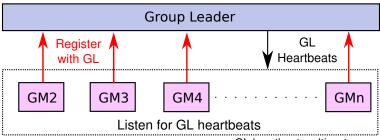


Group Manager Join

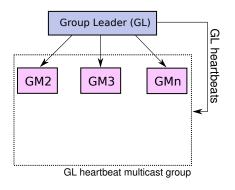


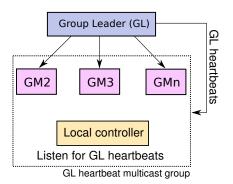
GL heartbeat multicast group

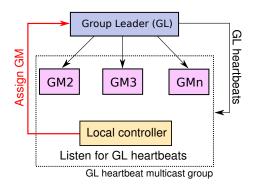
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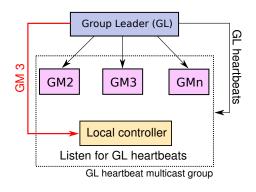


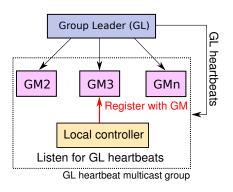
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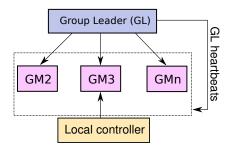






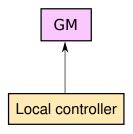


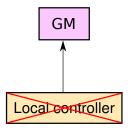


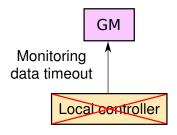


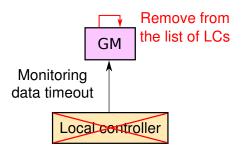
Hierarchy Reconstruction and Maintenance

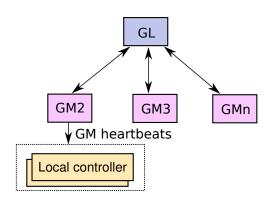
- Three kinds of failures
 - Local controller
 - Group manager
 - Group leader
- Two steps to tolerate failures
 - Error detection
 - 2 Recovery

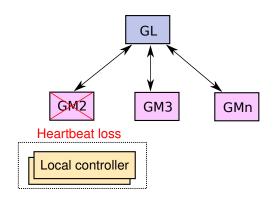


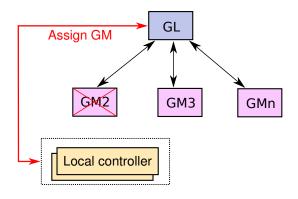


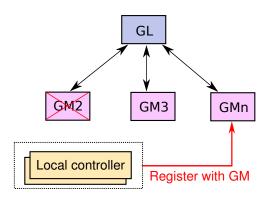


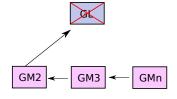


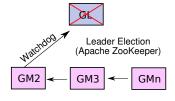


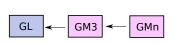


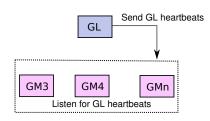




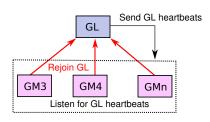












Hierarchy Management Evaluation

Scalability and self-healing

- Number of LC servers managed by a GM
- Number of GM servers managed by a GL
- Cost of the heartbeat mechanisms
- Cost of the self-healing mechanisms

Prototype implementation deployed on the Grid'5000 testbed



E. Feller, L. Rilling, and C. Morin. Snooze: A Scalable and Autonomic Virtual Machine Management Framework for Private Clouds. In the 12th IEEE/ACM International Symposium on Cluster, Cloud, and Grid Computing (CCGrid), May 2012.

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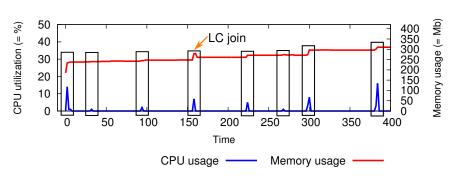
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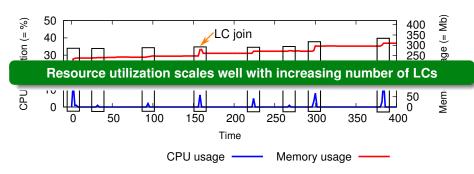
Group Manager Scalability

How does the GM server CPU and memory utilization scale with increasing number of LC servers?



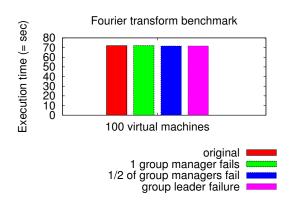
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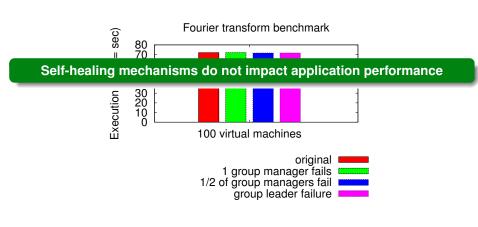
Cost of the Self-healing Mechanisms

What is the impact of the self-healing mechanisms on the application performance?



Cost of the Self-healing Mechanisms

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Outline



- Self-configuring and self-healing VM management system
- Self-optimizing integrated energy management approach

Mechanisms and Algorithms for Energy Efficiency

- How to favour idle times
 - Energy-efficient VM placement
 - Underload server detection and mitigation
 - Periodic VM consolidation
- Server overload detection and mitigation
- Power management
 - Automatic detection and power cycling of idle servers
 - Server wakeup when not enough resources are available

Approach	Algorithm	Resources	Placement	Underload, Overload Mitigation	VM Consoli- dation	Power manage- ment	Evaluation

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Entropy	Constraint programming	CPU, RAM	Yes (Consolidation)	Overload (Consolidation)	Yes	Server off/on	Real system and simulations

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Sandpiper	Greedy	CPU, RAM, Network	No	Overload	No	No	Real system
V-MAN	Greedy	Number of VMs	No	No	Yes	No	Simulation
Sercon	Greedy	CPU, RAM	No	No	Yes	No	Simulation
Borgetto et al. (2012)	Greedy	CPU, RAM	Yes	Underload, Overload	Yes	Server off/on	Simulation
VMware DRM	Greedy (Pro- prietary)	CPU, RAM	Yes	Underload, Overload	No	Server off/on	Real system

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Snooze	Greedy (ex- tensible)	CPU, RAM, Network	Yes	Underload, Overload	Yes (modi- fied Sercon)	Server off/on	Real system

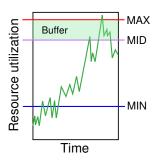
Underload and Overload Mechanism

- How to deal with underload and overload situations?
 - Detection of server underload/overload situations
 - Relocation of VMs from underloaded/overloaded servers

Underload and Overload Detection Approach

Local controllers periodically estimate their resource utilization based on locally aggregated VM resource utilization data

- Multi-dimensional
 - CPU
 - RAM
 - Network Rx
 - Network Tx

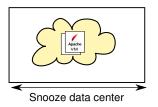


Underload Relocation Algorithm

Triggered by the GM in the event of server underload

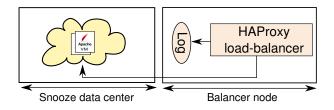
- Key ideas
 - Move VMs from underloaded LC to LCs with enough spare capacity
 - All-or-nothing approach: Either migrate all VMs or none
- Description
 - Sort VMs from underloaded LC in decreasing order of estimated utilization
 - Sort destination LCs in decreasing order of estimated utilization
 - Attempt to assign the VMs to the destination LCs starting from the first one
 - If some VM could not be assigned abort the algorithm
 - ... otherwise perform live migrations

Evaluation with an elastic web service



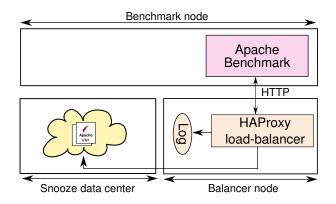
Deployed on 34 power-metered servers of the Grid'5000 testbed

Evaluation with an elastic web service



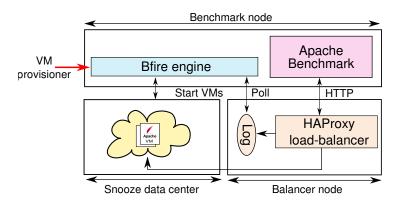
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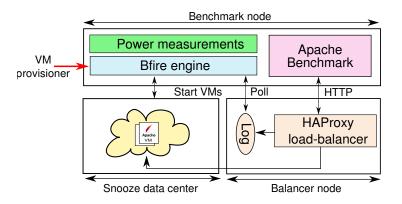
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Evaluation with an elastic web service



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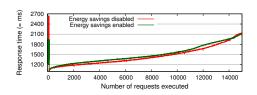
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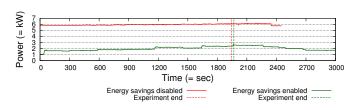
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Energy Saving Evaluation

Apache Benchmark Performance



Data Center Power Consumption



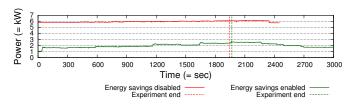
Energy Saving Evaluation

Apache Benchmark Performance



Limited performance degradation
Up to 67% energy savings for the evaluated application

Data Center Power Consumption



First Contribution Summary

- Self-configuring and healing hierarchical architecture
- Integrated energy management approach
 - VM placement and consolidation, server underload/overload mitigation, power management
 - Four-dimensional aggregation-based underload/overload mitigation
 - First implementation of the Sercon algorithm in a real system
- A robust prototype
- Experimentally validated on the Grid'5000 testbed

Second Contribution Presented

Virtual machine consolidation

Approach	Algorithms	Worst-case Complexity	Solution	Parallelization
Greedy	Sercon	Polynomial	Close to optimal	No
				

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Greedy	Sercon	Polynomial	Close to optimal	No
Mathematical programming	Constraint programming	Exponential	Optimal	Yes

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Mathematical programming	Constraint programming	Exponential	Optimal	Yes
Metaheuristics	Genetic algorithms, Ant Colony Optimization	Polynomial	Close to optimal	Yes

Approach	Algorithms	Worst-case Complexity	Solution	Parallelization
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Mathematical programming	Constraint programming	Exponential	Optimal	Yes
Metaheuristics	Genetic algorithms, Ant Colony Optimization	Polynomial	Close to optimal	Yes



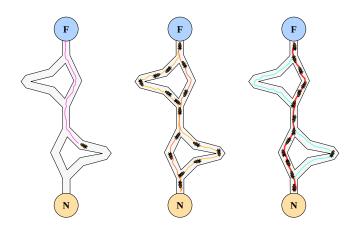
First attempt to apply Ant Colony Optimization on VM consolidation

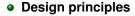
Ant Colony Optimization

- Ants work independently
- Indirect communication using pheromone in the environment
- Decisions are taken probabilistically

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- Ants compute solutions concurrently
- Best solution is preserved
- Pheromone on VM-server pairs
- Probabilistic pair choice





Server 2



Design principles

Ants compute solutions concurrently

- Best solution is preserved
- Pheromone on VM-server pairs
- Probabilistic pair choice

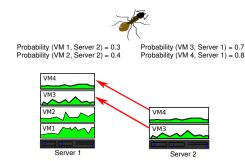
Probability (VM 1, Server 2) = 0.3 Probability (VM 2, Server 2) = 0.4 Probability (VM 3, Server 1) = 0.7Probability (VM 4, Server 1) = 0.8





Design principles

- Ants compute solutions concurrently
- Best solution is preserved
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Design principles

- Ants compute solutions concurrently
- Best solution is preserved
- Pheromone on VM-server pairs
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Server 1



Algorithm components

- Objective function
- Probabilistic pair selection rule
- Pair pheromone update rule

VM Consolidation Scalability Issues

- VM consolidation by nature is not scalable
 - Computing optimal solutions is exponential in time and space
 - Solution quality degrades at scale

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Desirable properties

- Scalability with increasing number of servers and VMs
- High packing efficiency (PE)

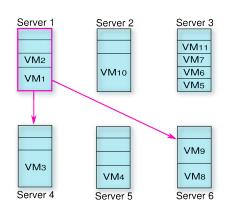
$$PE := \frac{\text{Number of released servers}}{\text{Total number of servers}} \times 100$$

Minimize the number of migrations

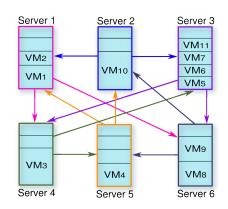
Fully Decentralized VM Consolidation System

- Servers maintain only a partial system view
- VM consolidation is applied within these partial views
- Partial views are modified periodically and randomly

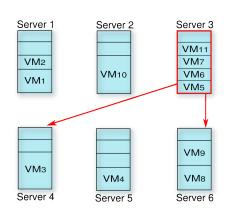
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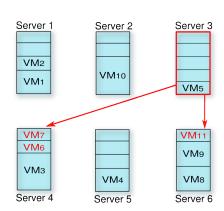
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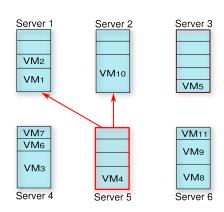
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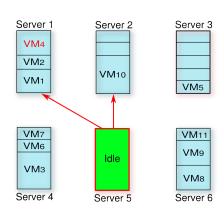
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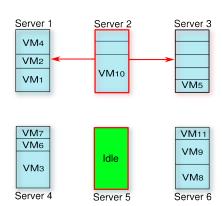
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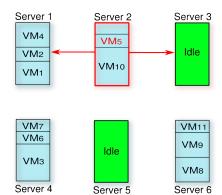
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Fully Decentralized VM Consolidation System Evaluation

Criteria

- Scalability
- Packing efficiency
- Number of migrations

Experiments

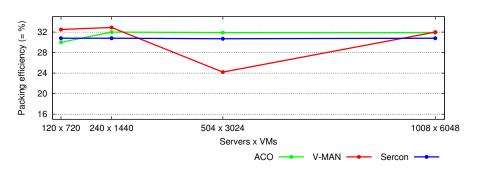
- Comparison of different VM consolidation algorithms
 - Sercon
 - V-MAN
 - Our ACO-based VM consolidation algorithm
- Comparison with a centralized system

Evaluated by emulation

E. Feller, C. Morin, and A. Esnault. A Case for Fully Decentralized Dynamic VM Consolidation in Clouds. In the 4th IEEE International Conference on Cloud Computing Technology and Science (CloudCom) (Best Paper Finalist), December 2012.

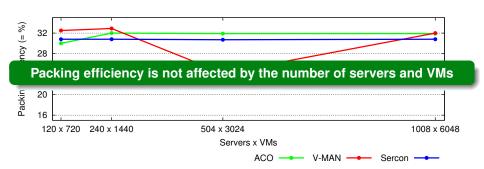
Scalability

How does the system scale in terms of its packing efficiency with increasing number of servers and VMs?



Scalability

How does the system scale in terms of its packing efficiency with increasing number of servers and VMs?



Comparison With a Centralized System Topology

What is the packing efficiency and number of migrations compared to a centralized system?

Topology	Algorithm	Migrations	Packing Efficiency (%)
Centralized	Sercon	1920	31.7
Centralized	ACO	Failed	Failed
P2P	V-MAN	4189	32.0
P2P	ACO	4015	31.9
	Sercon	1872	30.8

Experiments with 1008 servers and 6048 VMs

Comparison With a Centralized System Topology

Algorithm

What is the packing efficiency and number of migrations compared to a centralized system?

	lobology	Algorialin	iviigrations	Tacking Efficiency (70)	
	Centralized	Sercon	1920	31.7	
	Oontraiizoa	۸۲۸	Failad	Failad	
Packing efficiency and number of migrations close to a centralized system					d system
		ACU	4015	১ ৷ . ৬	
		Sercon	1872	30.8	

Migrationa Booking Efficiency (9/)

Experiments with 1008 servers and 6048 VMs

Second Contribution Summary

- ACO-based VM consolidation algorithm
- Fully decentralized VM consolidation system
- Validated on the Grid'5000 experimentation testbed
 - Scalable with increasing numbers of servers and VMs
 - Packing efficiency close to a centralized system

Criteria	Best algorithm	2nd	3rd
#Migrations	Sercon	ACO	V-MAN
Packing efficiency	V-MAN	ACO	Sercon

Conclusion

Snooze: autonomic and energy-efficient VM management system for large-scale laaS clouds

- Self-configuring and healing hierarchical architecture
- Platform to evaluate VM management algorithms in a real system
- Open-source software (http://snooze.inria.fr)
 - External users: IRIT Toulouse, EDF R&D, LIFL, LBNL, and Medion Seattle
 - Support: Inria technological action

Algorithms for energy efficiency

- Evaluation of an integrated approach
 - First implementation of Sercon consolidation algorithm in a real system
 - Novel approach for underload/overload management
 - Up to 64% of energy savings
- First ACO-based placement and consolidation algorithms
 - Viable approach in a fully decentralized system

Short-term Perspectives

- Further evaluate the Snooze system
 - Larger-scale experiments
 - Real-world workloads
 - Hierarchy energy overheads
- Exploit Snooze to experimentally compare state of the art VM management algorithms
- Further increase the Snooze hierarchy autonomy and energy-efficiency
 - Re-balance the hierarchy dynamically
 - Remove local controller/group manager distinction
 - Power-cycle idle GMs

Long-term Perspectives

- Metrics for better capturing aggregated resource utilization data
- Improving consolidation
 - Co-location and anti-colocation constraints
 - Consider VM resource demand complementarities
 - Data center network topology aware consolidation
 - Consolidation interval predictions
- Thermal management

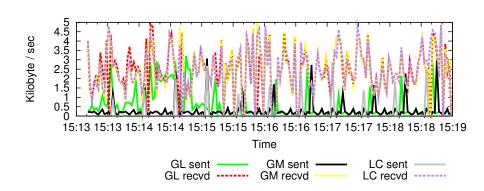
Questions?

Thank you for your attention!

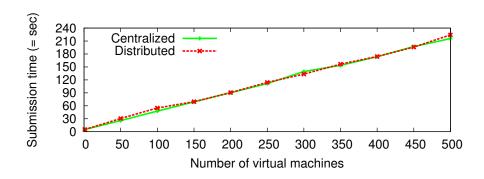
Discussion

Backup slides

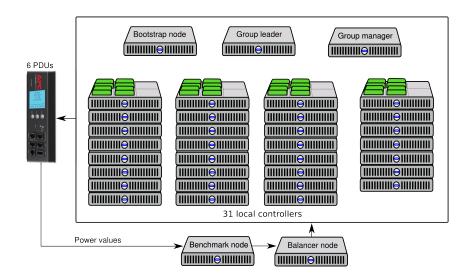
Snooze Heartbeat Overhead



Snooze Submission Time



Energy Management Data Center



Energy Management Parameters

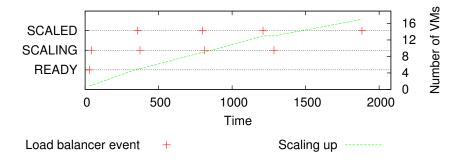
Resource	MIN, MID, MAX
CPU,	0.2, 0.9, 1
Memory	0.2, 0.9, 1
Network	0.2, 0.9, 1

Policy	Algorithm
Dispatching	RoundRobin
Placement	FirstFit
Overload	Greedy
Underload	Greedy
Consolidation	Sercon

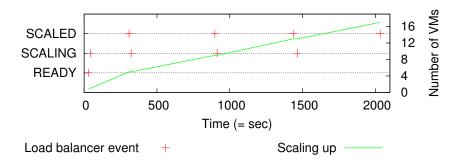
Parameter	Value
Packing density	0.9
Monitoring backlog	15
Resource estimators	average
Consolidation interval	10 min

Parameter	Value	
Idle time threshold	2 min	
Wakeup threshold	3 min	
Power saving action	shutdown	
Shutdown driver	system	
Wakeup driver	IPMI	

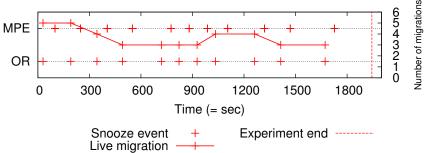
Bfire Events With Energy Savings Disabled



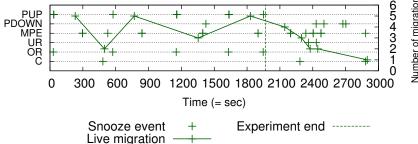
Bfire Events With Energy Savings Enabled



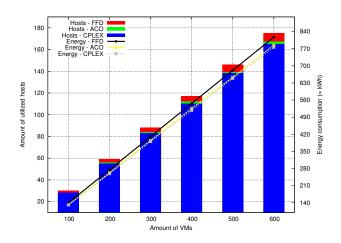
Snooze Events No Energy Savings Disabled



Snooze Events With Energy Savings Enabled



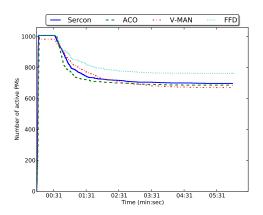
VM Placement Evaluation



Fully Decentralized VM Consolidation - Emulator Parameters

Parameter	Value	
Number of PMs and VMs	1008 (resp. 6048)	
Experiment duration	360s	
Consolidation interval	30s	
Shuffling interval	10s	
Neighbourhood size	16 PMs	
Considered resources	CPU, memory and net-	
	work	
PM total capacity vector	(48, 26, 20)	
VM requested capacity vectors	(0.2, 0.5, 0.1), (1, 1, 1), (2,	
	1, 1), (4, 2, 2), (8, 4, 4),	
	(16, 8, 4)	

Fully Decentralized VM Consolidation - Number of Active Servers



Fully Decentralized VM Consolidation - Number of Migrations

